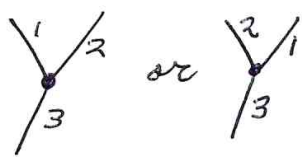


Penrose Formula, Perfect Matching  
Polynomials, Virtual Knot Theory  
and Khovanov Homology

Louis H Kauffman  
UC & NSU

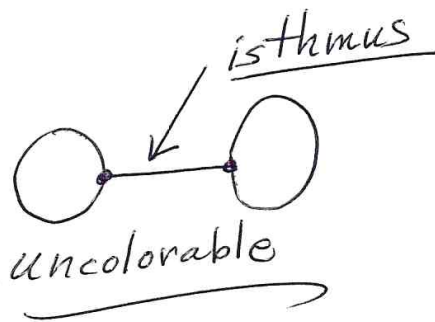
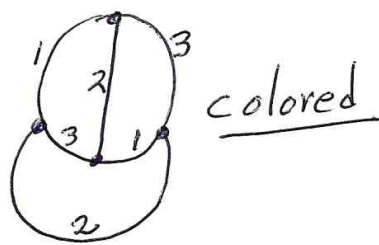
(joint work with Scott Baldridge,  
and William Rushworth)

1. Coloring Trivalent Graphs



Rule: 3 distinct  
edge colors at  
each node.

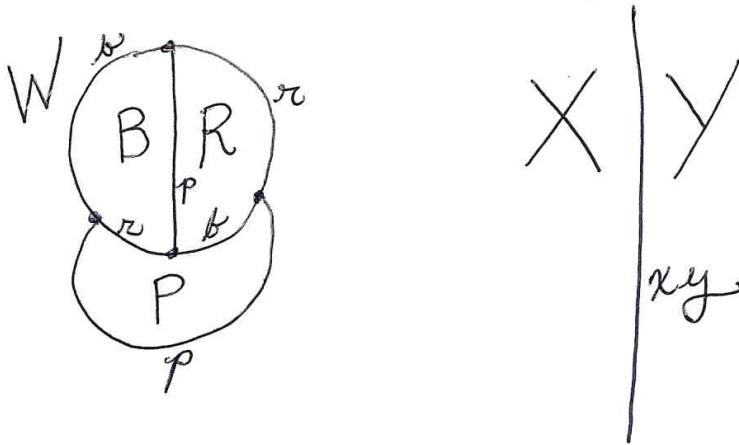
example



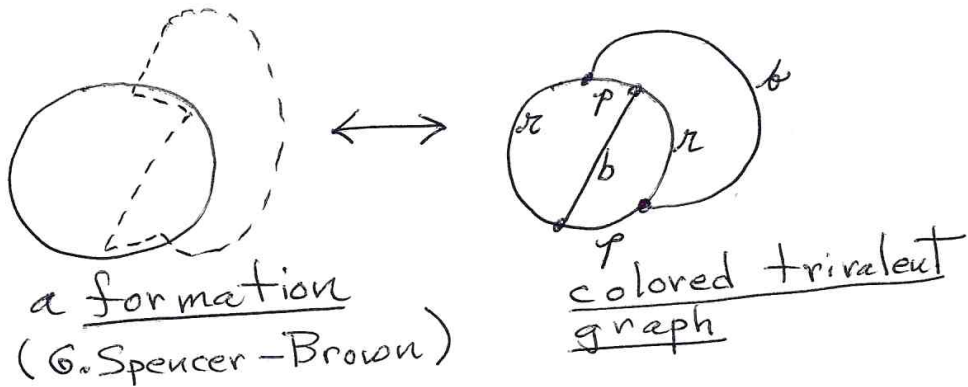
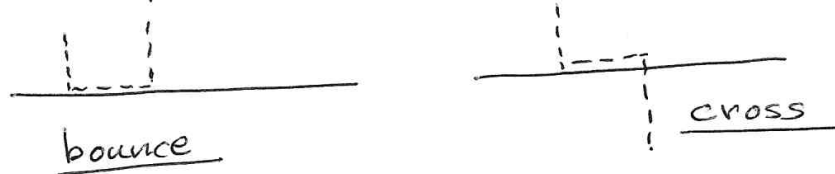
Theorem (Peter Guthrie Tait). 3-Colorability  
of isthmus-free plane trivalent  
graphs is equivalent to the  
Four Color Theorem (4CT).

example

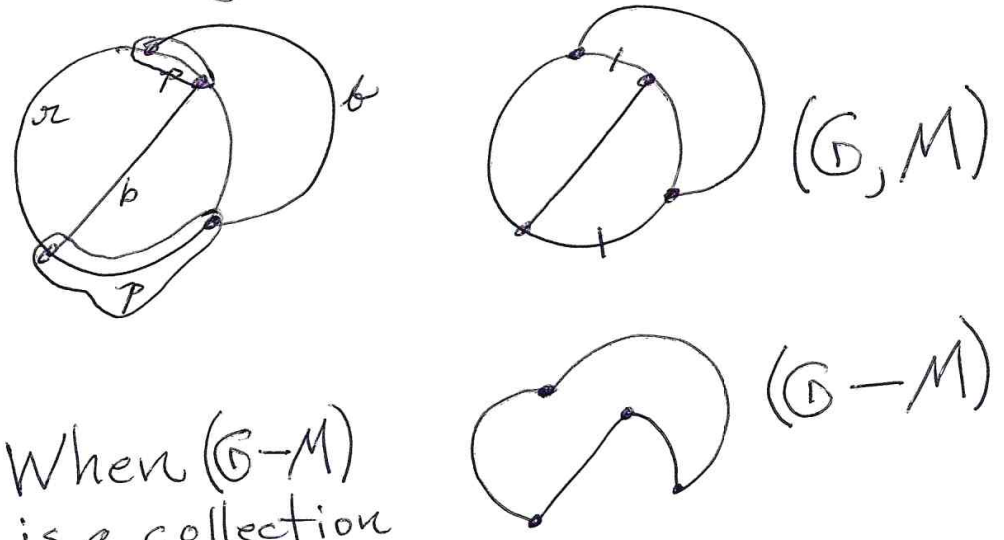
(a) Let colors be  $r$  (red),  $b$  (blue) and  $p$  (purple). Take four colors  $\{W, R, B, P\} \cong \mathbb{Z}_2 \times \mathbb{Z}_2$  with  
 $W = \text{Identity}, R^2 = B^2 = P^2 = W$   
 $RB = P, BP = R, PR = B.$



(b)         $r$             $b$             $p$   
 Let red and blue curves interact.



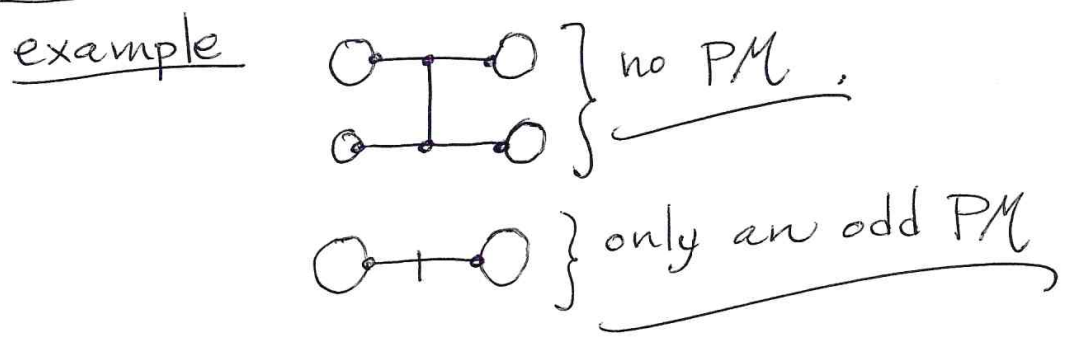
Note that every colored trivalent graph selects an even perfect matching (choose the purple edges).



When  $(G - M)$  is a collection of even cycles, we say  $M$  is an even perfect matching.

$G$ is 3-colorable	$\iff$	$G$ has an <u>even</u> PM.
--------------------	--------	----------------------------

Easier Than 4CT: Every isthmus-free trivalent graph has a PM.



2. Penrose Formula




$$[\text{X}] = [\text{Y}] - [\text{Z}]$$

$$[0] = 3$$




$\Rightarrow [G] = \# \text{ of } 3\text{-colorings of } G$   
 when  $G \hookrightarrow \mathbb{R}^2$ .

(From: R. Penrose, On Applications of Negative Dimensional Tensors, in "Combinatorial Mathematics and Its Applications", ed by D.J. Welsh (Academic Press 1971) )

example

(a)   $\rightarrow$    $-$    $= 3^2 - 3 = 6$

$$[\text{circle with line}] = 6.$$

(b)   $\rightarrow$    $-$    $= 3 - 3 = \emptyset$

$$[\text{two circles connected}] = \emptyset.$$

(c)  $[\text{circle with line and dot}] = [\text{circle with line}] - [\text{figure-eight with dot}] = 3[\text{vertical line}] - [\text{vertical line}] = 2[\text{vertical line}]$

(d)  $\rightarrow$   $-$   $+ \dots$

$$= \emptyset - \dots + \dots$$

$$= -3 + 00$$

$$= -3 + 3^2$$

$\left[ \text{Diagram} \right] = 6.$

(e)  $\rightarrow$   $-$   $\rightarrow \emptyset$

(f)  $\rightarrow$   $\rightarrow$   $-$   $\rightarrow 3 - 3^2 = -6$

*not a plane diagram*

(g)  $\rightarrow$   $-$   $\rightarrow$   $-$   $= \emptyset.$

$K_{3,3}$   
non-planar

Note:  $K_{3,3}$  has 12 distinct colorings.



# Proof of Penrose Formula

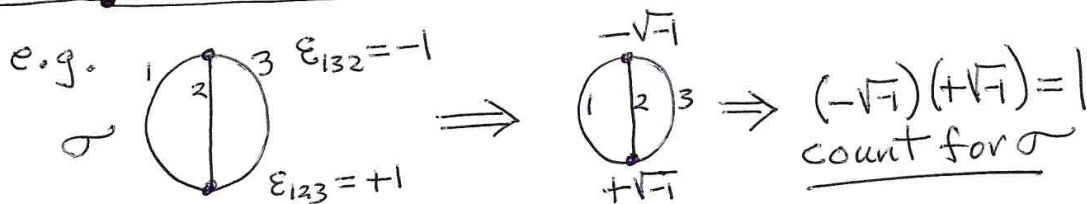
$$[\textcircled{1}] \stackrel{\text{def}}{=} \sum_{\sigma \in \text{Colorings}(\textcircled{1})} \prod (\pm \sqrt{-1}) = \mathcal{C}(\sigma)$$

And

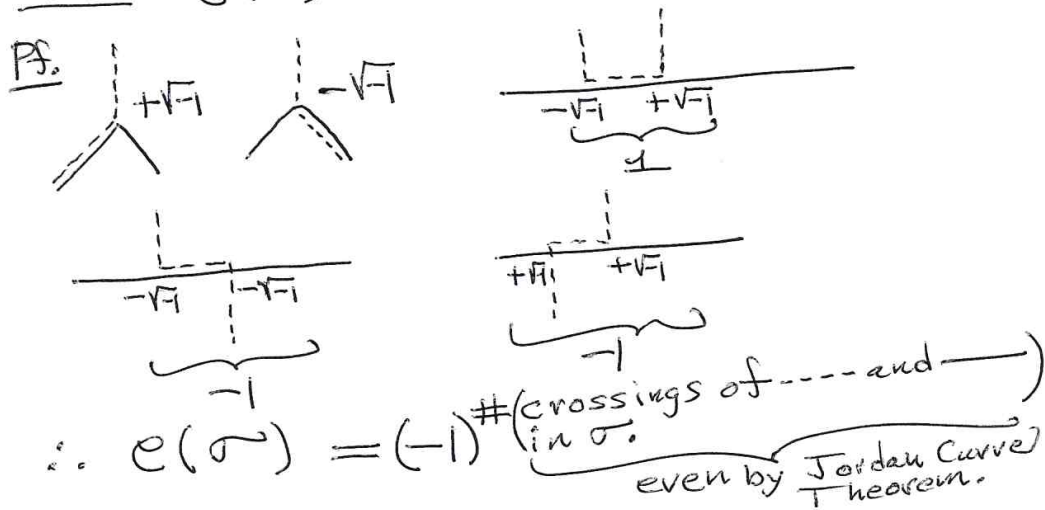
Epsilon Identity

$$\pm \sqrt{-1} = \sqrt{-1} \begin{array}{c} a & b \\ \diagdown & / \\ & c \end{array} = \sqrt{-1} \epsilon_{abc}$$

$\epsilon_{abc} :$	$\epsilon_{123} = +1$	}	$\epsilon_{123} = \epsilon_{231} = \epsilon_{312}$
	$\epsilon_{213} = -1$		$\epsilon_{213} = \epsilon_{132} = \epsilon_{321}$
	$\epsilon_{113} = 0$		



Claim:  $\mathcal{C}(\sigma) = 1$  for each  $\sigma$ .



$$\mathcal{C}(\sigma) = 1 //$$

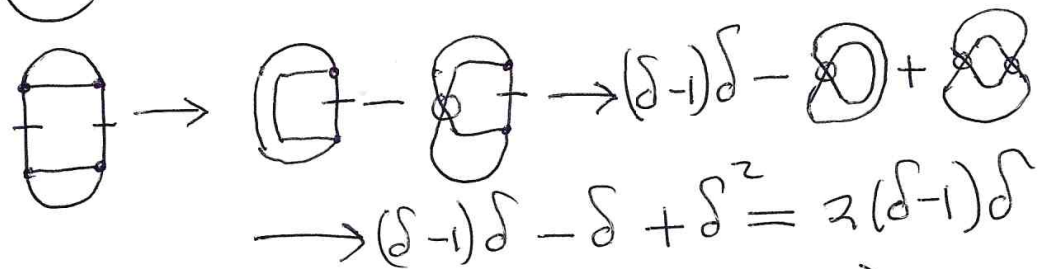
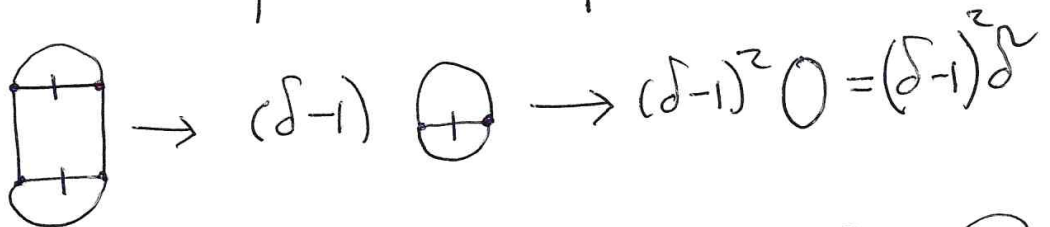
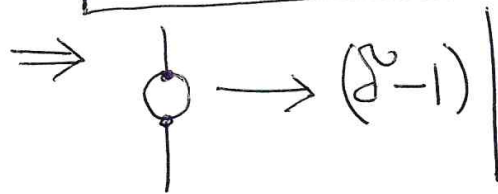
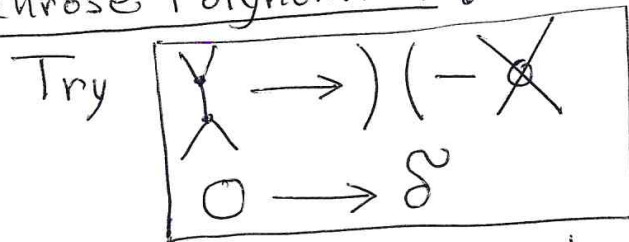
$$-\sum_i \epsilon_{abi} \epsilon_{icd} = \delta_a^a \delta_c^b - \delta_c^a \delta_d^b$$

(The fact that  $[K_{3,3}] = \emptyset$  proves that  $K_{3,3}$  is not planar.)

⑥

Penrose formula can be modified to count colorings for non-planar graphs. See (LHK, *ILJ, Math. Vol. 60 No. 4, (2016) PP 251-271*).

3. Penrose Polynomial?



While  $(3-1)^2 \cdot 3 = 2 \cdot (3-1) \cdot 3,$

$(\delta-1)^2 \delta \neq 2(\delta-1)\delta$

and  $\Leftrightarrow \delta-1=2 \Leftrightarrow \delta=3.$

Such an extension produces a perfect matching polynomial.

# 4° Virtuality of Graphs, Knots, Links

⑦

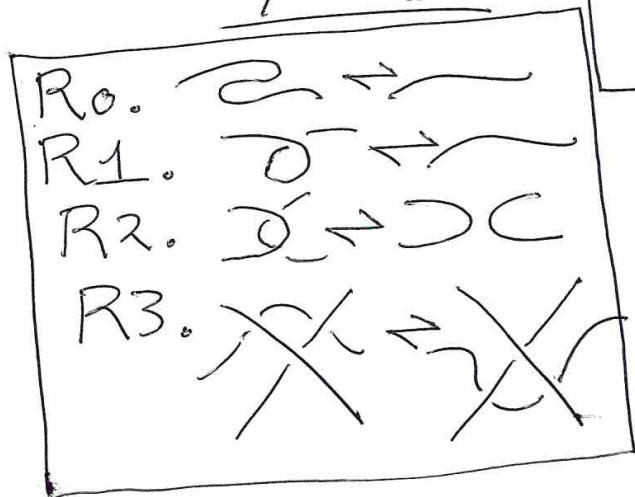
Remark. All our calculations involving ~~use~~ virtual graph equivalence

Virtual Knot Theory  
Virtual Knot, Link Diagrams

with  
Reidemeister Moves

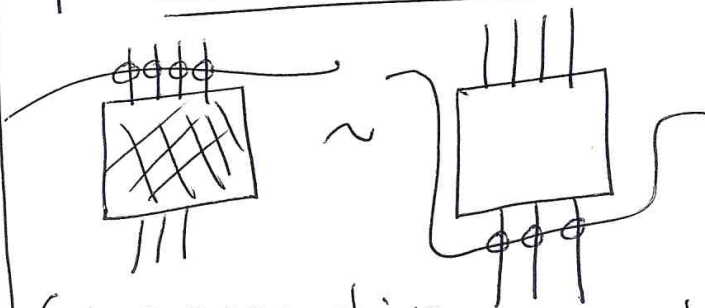
+  
Detour Moves

Reidemeister Moves

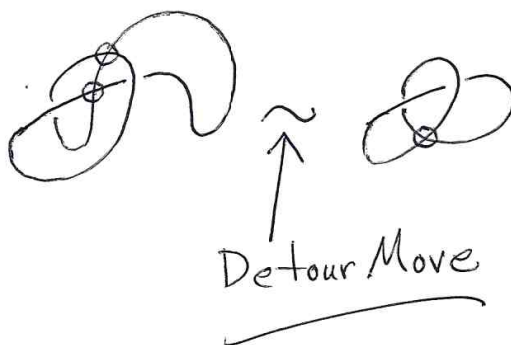


planar isotopy  
 (preserving cyclic orders at graphical vertices)

+ detour moves



(a consecutive sequence of virtual crossings can be deleted and replaced elsewhere by another consecutive sequence of virtual crossings).





Let  $(\mathbb{G}, M)$  be a (virtual) plane diagrammatic trivalent graph with given perfect matching  $M$ .

Define a 3-variable perfect matching polynomial  $[\mathbb{G}, M]$  via:

$$[\text{Y}] = A[\text{)}(\text{]} + B[\text{X}]$$

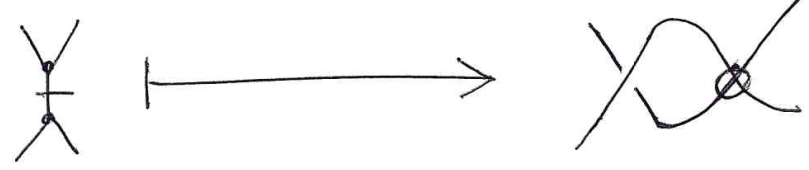
$$[\text{O}] = \delta$$

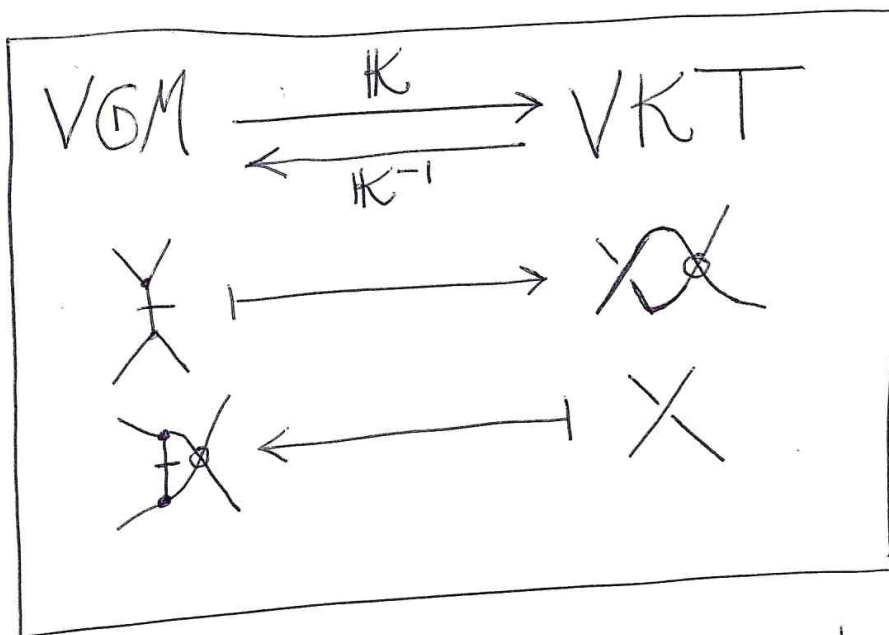
This can be used to discriminate many different perfect matchings on  $\mathbb{G}$ .

VGM = virtual graphs with perfect matching.

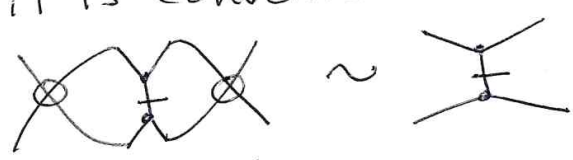
VKT = virtual knot and link diagrams.

$$\text{VGM} \xleftrightarrow{\mathbb{K}_G} \text{VKT}$$

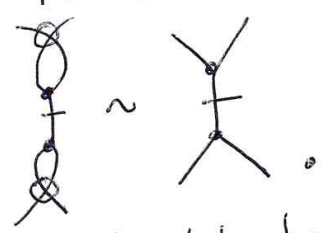




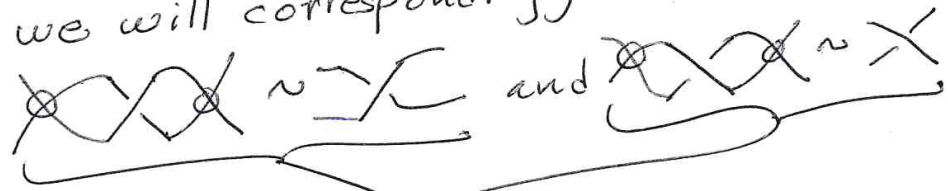
For some purposes, as we shall see, it is convenient to assume that



and that

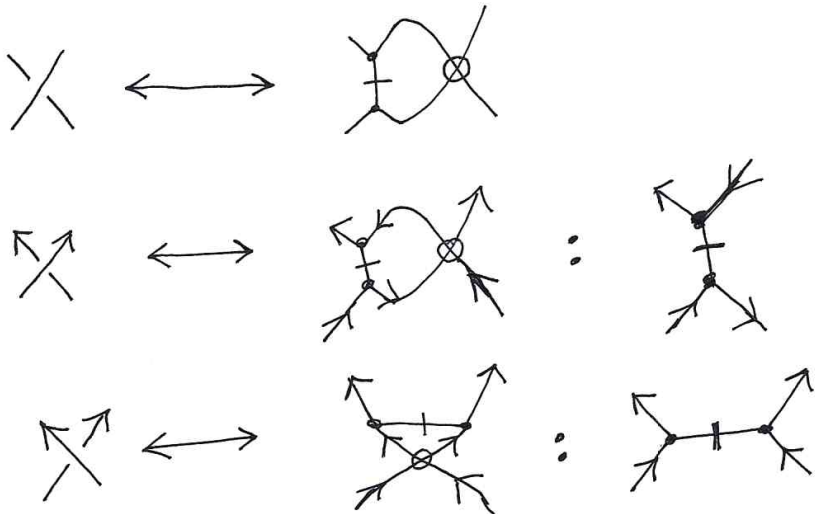


I will do that here, but take note that we will correspondingly demand that



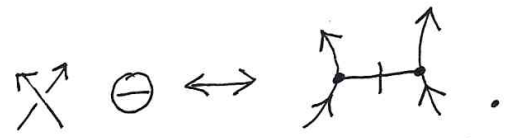
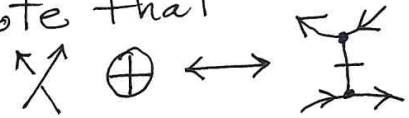
Z-equivalence

Orientation

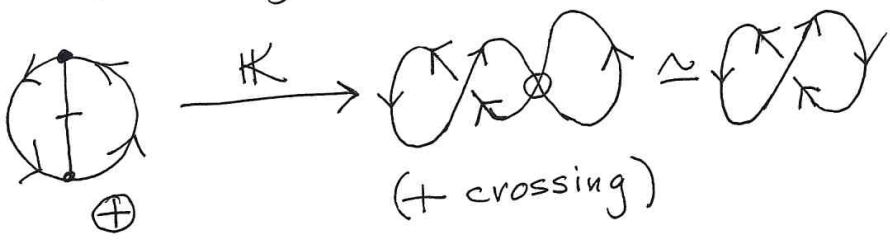


These two cases comprise the allowed orientations of the edges in  $\mathbb{G}-M$  for  $(\mathbb{G}, M)$  a perfect matching.

Note that



(ignoring the requisite virtual crossings)



9

Define for  $K \in VKT$ ,

$$[K] = [K^{-1}(K)]$$

So:  $[X] = [X \otimes]$

$$= A[\text{)(}] + B[X \otimes]$$

$$[X] = A[\text{)(}] + B[\text{)](}]$$

and  $[O] = d$ .

Thus the 3-variable PM poly on trivalent graphs transfers to the well-known 3-variable bracket on virtual link diagrams.

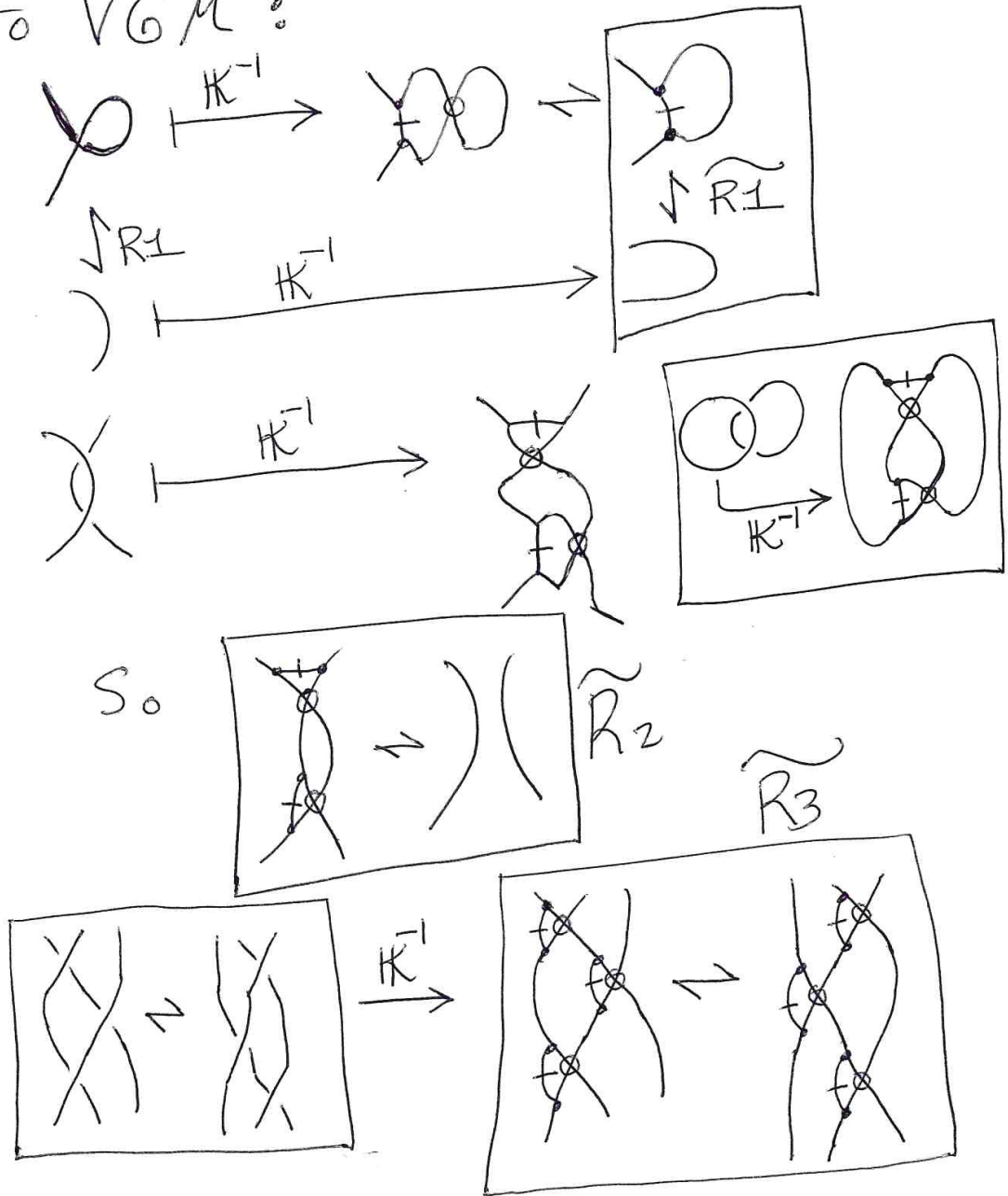
Taking  $B = A^{-1}$  and  $d = -A^2 - A^{-2}$ , we then have a bracket polynomial for perfect matchings

$$\langle X \rangle = A \langle \text{)(} \rangle + A^{-1} \langle X \otimes \rangle$$

$$\langle O \rangle = -A^2 - A^{-2}$$

(10)

and correspondingly a Jones polynomial that is invariant under all 3 R moves. Perfect matching polynomials have topological properties & we can transfer the R-moves over to VGM:





At  $A=i$  we have

$$\langle \text{X} \rangle = i() (-\text{X})$$

$$\langle \text{O} \rangle = 2$$

and  $\langle \text{Y} \rangle = i() (-\text{Y})$

$$\langle \text{O} \rangle = 2.$$

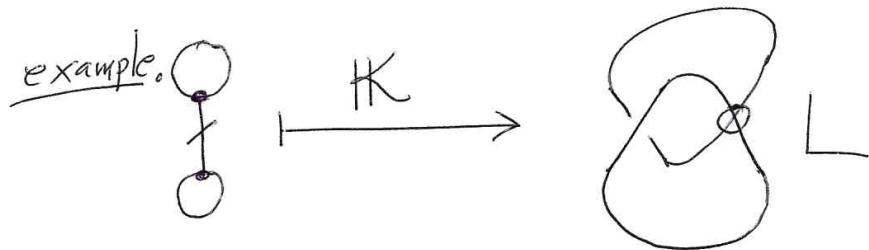
Transferring a result of Scott Baldridge about the perfect matching polynomial, we find

arXiv: 1812.10346 and 1810.07302

$$\langle K \rangle \neq 0 \iff \text{every link component of } K \text{ meets } K \text{ in an even number of virtual crossings}$$

$$\langle \textcircled{D}_M \rangle \neq 0 \iff \text{the perfect matching } M \text{ on } \textcircled{D} \text{ is even.}$$

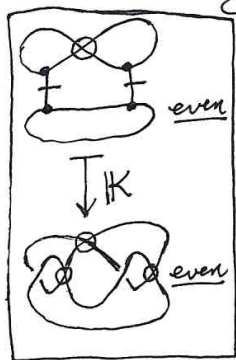
When  $(\textcircled{D}, M)$  is an even PM then  $|\langle \textcircled{D}, M \rangle| = 2^{\# \text{cycles}(\textcircled{D}-M)}$



Compute  $\langle L \rangle$  at  $A=i$ :

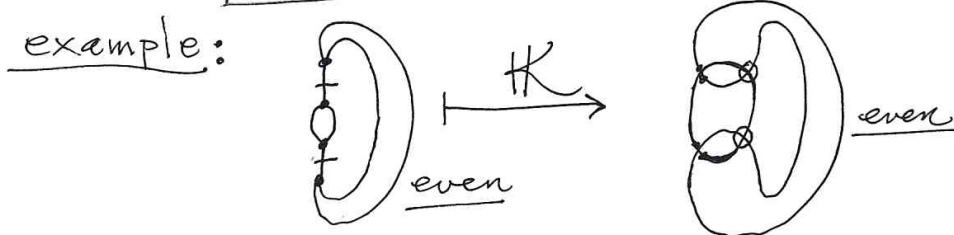
$$\langle L \rangle = i \langle \text{link with 2 crossings} \rangle - i \langle \text{link with 2 crossings} \rangle = \emptyset.$$

example. Call a virtual link even if any component shares <sup>(with the other components)</sup> an even number of virtual crossings. Then  $\mathbb{K}^+$  (any virtual knot diagram) and  $\mathbb{K}^-$  (any <sup>even</sup> virtual link diagram) are even perfect matchings. Thus we produce many already colorable graphs.



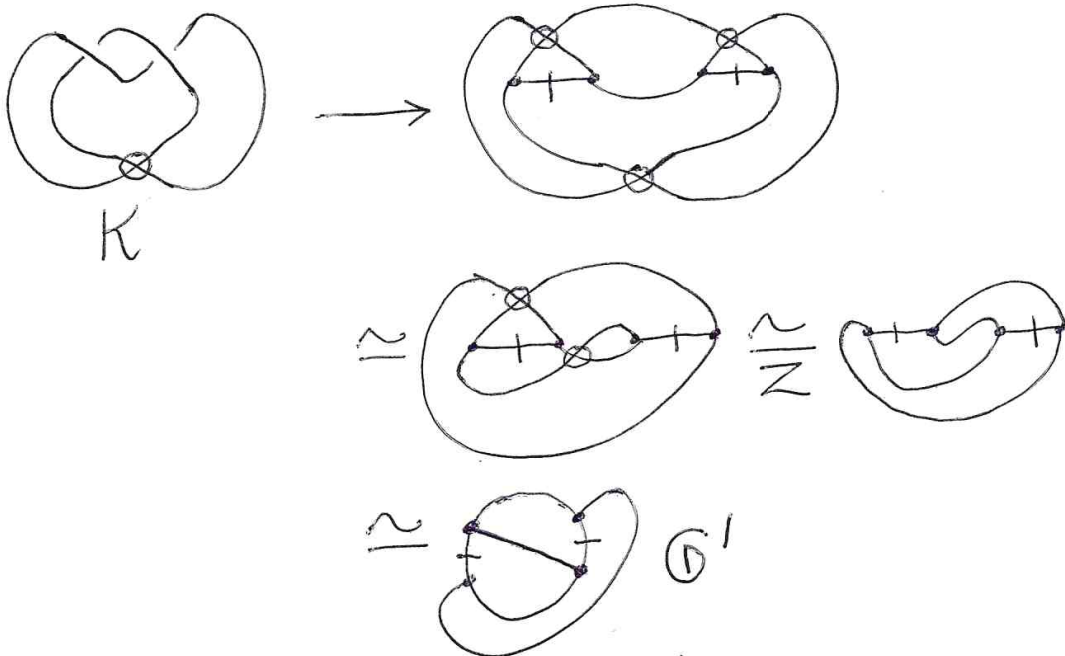
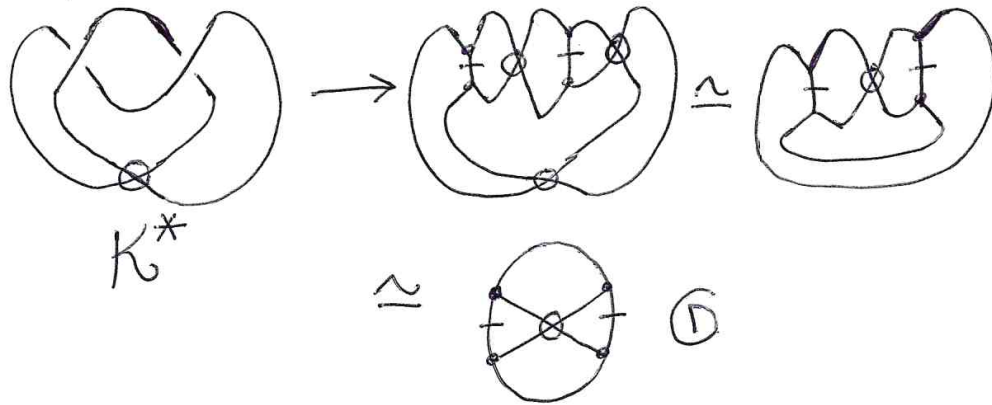
The 4CT says that, among them, are all isthmus-free plane graphs.

Will this topological framework yield deeper insight into the 4CT?



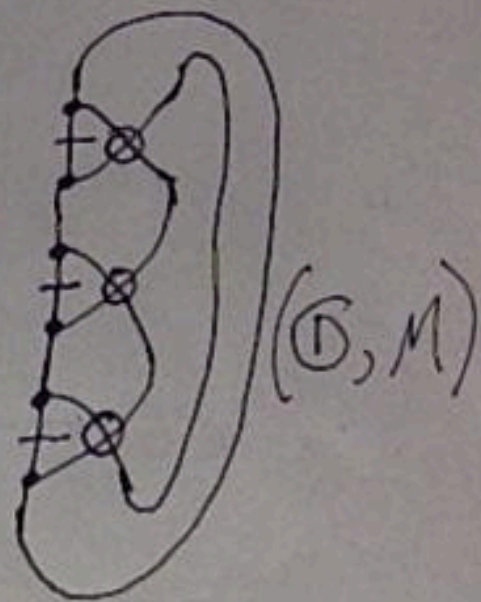
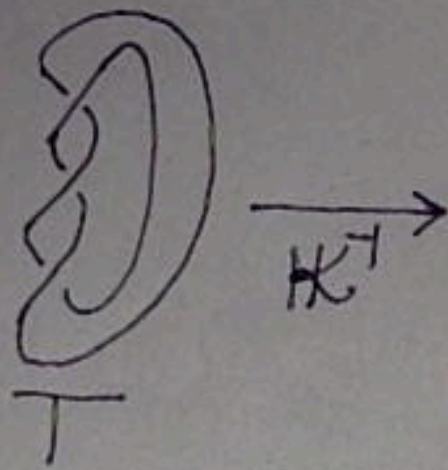
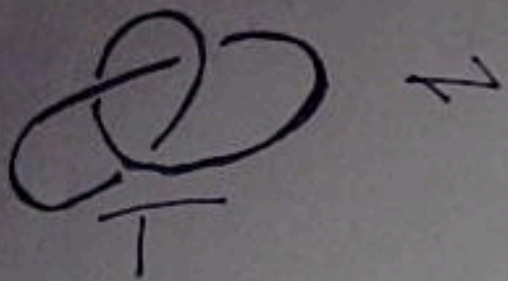
example

(13)

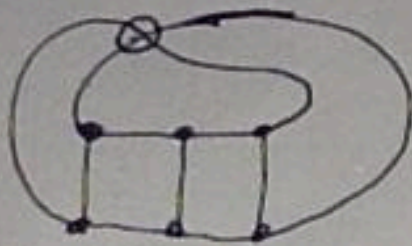


Since  $\langle K \rangle = A^2 + 1 - A^{-4}$   
 $f_K = A^{-4} + A^{-6} - A^{-10}$   
 $f_{K^*} \neq f_K$   
 $\Rightarrow K^* \not\cong K$   
 $\Rightarrow G$  and  $G'$  are distinct perfect matching graphs.

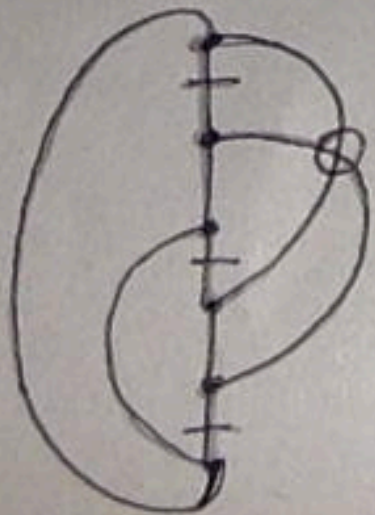
Note that both  $G$  and  $G'$  are even perfect matchings.



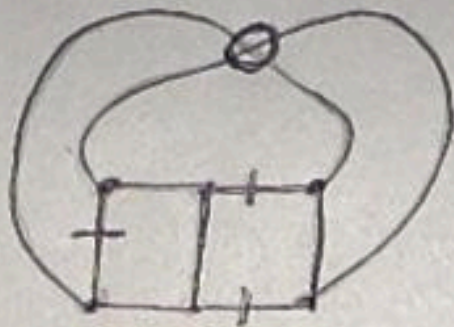
①  $\frac{\sim}{\text{graph}} K_{3,3}$



② 12

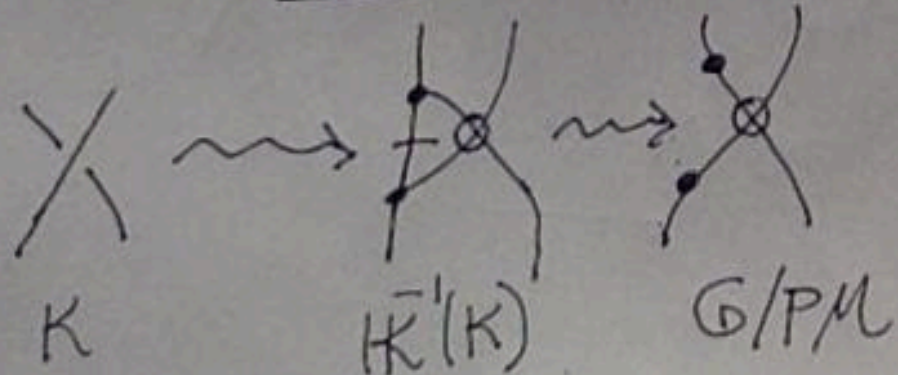


12





Note that  $K^{-1}(K) - \{\text{PM edges}\} = \mathbb{G}/\text{PM}$   
 $\equiv \bar{K} = \underline{K \text{ with all crossings made virtual.}}$



Theorem.  $K$  a classical knot or link diagram, then  $K^{-1}(K) = (\mathbb{G}, \mathcal{M})$

is an even perfect matching.  
Hence  $\mathbb{G}$  is 3-colorable.

Proof. By the Jordan Curve theorem any component in  $\mathbb{G}/\text{PM}$  will have an even number of virtual crossings.  
 $\therefore$  an even number of dots as above.  
This proves that  $(\mathbb{G}, \mathcal{M})$  is even. //



Proposition. Given a virtual trivalent graph  $\mathcal{G}$  with a PM and a single circuit in  $(\mathcal{G}-M)$ , then  $(\mathcal{G}, M)$  is an even PM, and hence  $\mathcal{G}$  is 3-colorable.

Proof.  $e = e' + e_{PM}$  where  
 $e' = \# \text{edges}(\mathcal{G}-M)$ .  $e_{PM} = \# \text{of PM edges}$ .

$$3v = 2e, \quad 2e_{PM} = v.$$

$$\therefore \frac{3v}{2} = e' + \frac{v}{2} \Rightarrow v = e'$$

$$\text{and } 3v = 2e \Rightarrow v \text{ even.}$$

$$\therefore e' \text{ is even.} \quad //$$

Corollary. Let  $K$  be any virtual knot diagram (one component), then  $\mathcal{G}(K)$  ( $\mathcal{K}(K) = (\mathcal{G}(K), M)$ ) is a 3-colorable trivalent graph.

Proving Scott Baldridge's result via the virtual link category.

Let  $[><] = [<>] - [>>]$

$[O] = 2.$

Use oriented diagram with  $wc = n_+ - n_-$ ,  $n_+ = \#(\nearrow)$ ,  $n_- = \#(\searrow)$

Define  $J_K = (-1)^{n_-} [K] (\Rightarrow J_{\nearrow} = J_{\searrow})$

Then  $J_K$  is invariant under all RMs.

(e.g.  $J_{\searrow} \rightarrow = J_{\searrow} - J_{\searrow} = 2J_{\searrow} - J_{\searrow} = J_{\searrow}$ ,  
 $J_{\searrow} \rightarrow = -(J_{\searrow} - J_{\searrow}) = -(J_{\searrow} - 2J_{\searrow}) = J_{\searrow}$ .)

Theorem.  $J_K = 2^{|K|}$  if  $K$  is even and  $|K| = \#$  components of  $K$ .  $J_K = \emptyset$  if  $K$  is odd.  
(See previous defn of even and odd.)

Proof. The proof is by induction on the number of classical crossings in the diagram  $K$ . First suppose that  $K_i$  is a component of  $K$  and that  $\nearrow = c$  is a crossing between  $K_i$  and the rest of  $K$ . We can assume that all other crossings in  $K$  are positive (since  $J$  is unchanged by crossing switch).

Then  $J_{\nearrow} = J_{\searrow} - J_{\searrow}$  and each link on the right has one less component, with  $K_i$  now welded to a component  $K_j$  of  $K$ . In the second case we must re-orient the diagram to  $\searrow$ .

If the  $K_i$  contribution  $\lambda(K_i) = 1 + k$  is the count of classical intersections of other components with  $K_i$ , then  $\lambda \rightarrow k$  as contribution to  $\searrow$  and  $\lambda \rightarrow -k$  as contribution to  $\searrow$ . Thus  $\searrow$  acquires  $k$  negative crossings and we have

$J_K = J_{\searrow} - (-1)^k J_{\searrow}$ .  
Note that  $(-1)^{k+1} = (-1)^{\pi(K_i)}$  where  $\pi(K_i) = \#$  of virtual crossings between  $K_i$  and rest of  $K$ .

Thus, by induction

$$J_{\rightarrow} = 2^{|\rightarrow|} \text{ or } \phi = 2^{|\rightarrow|} \cdot \epsilon$$

$$J_{\rightarrow\leftarrow} = 2^{|\rightarrow\leftarrow|} \text{ or } \phi = 2^{|\rightarrow\leftarrow|} \cdot \epsilon'$$

and  $|\rightarrow| = |\rightarrow\leftarrow| = |\rightarrow| - 1$

and  $\epsilon' = \epsilon = \phi$  or  $\perp$ .

Thus  $J_K = (1 + (-1)^{\pi(K_i)}) 2^{|\rightarrow|} \cdot \epsilon$ .

The same argument applies when both lines of  $\rightarrow\leftarrow$  are on the same component (so that  $\rightarrow$  and  $\rightarrow\leftarrow$  have extra components). The theorem follows by induction since

$$(1 + (-1)^{\pi(K_i)}) = \begin{cases} 2 & \text{when } \pi(K_i) \text{ is even,} \\ 0 & \text{when } \pi(K_i) \text{ is odd.} \end{cases}$$

Remark. Generalize to the following form of the Jones polynomial.

$$[\times] = [\curvearrowright] - q[\curvearrowleft]$$

$$[O] = q + q^{-1}$$

$$J_K = (-1)^n q^{n-2n-} [K]$$

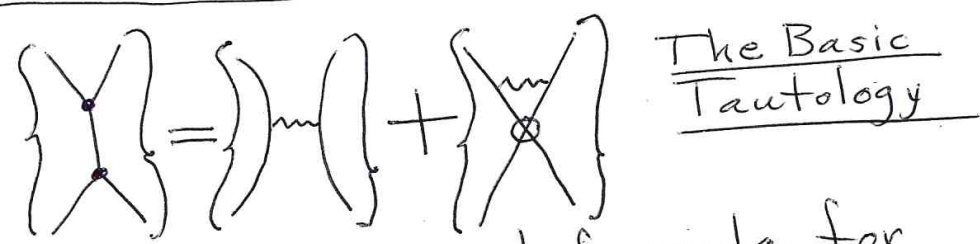
This version is useful for keeping track of grading in Khovanov Homology, and corresponds (by  $\mathbb{K}^{-1}$ ) to the perfect matching polynomial

$$P_Y = P_{\text{)}(-q P_{\text{X}}$$

$$P_O = q + q^{-1}$$

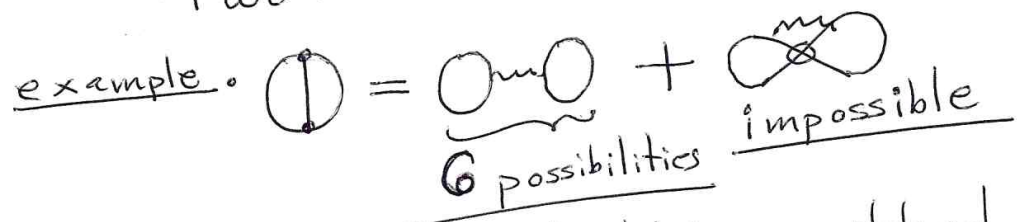


# 5.0 Return to Coloring



This is a general formula for counting 3-colorings. Here

$\overbrace{)}^m$  means that the two arcs are colored differently.



We can implement this method by choosing a perfect matching and expanding on its edges.  
The perfect matching need not be even.

example

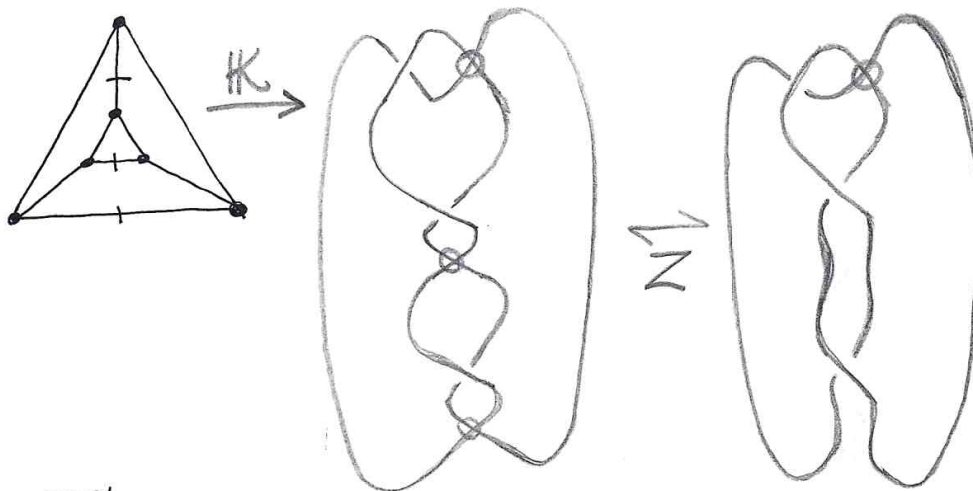
We start with any PM and locate all even PM's!

$\textcircled{1}$   $\vdots$  The "Astate" of  $\textcircled{1}$  is colorable with 3 colors.

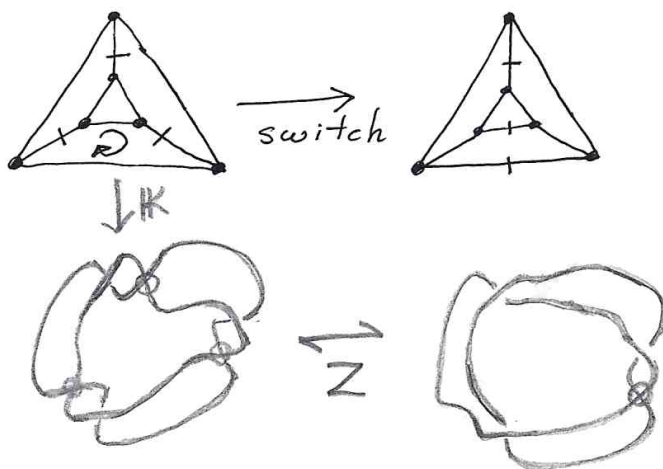
Take purple PM:

even

Whenever a graph is colorable, it has an even PM, and a corresponding even virtual link diagram.

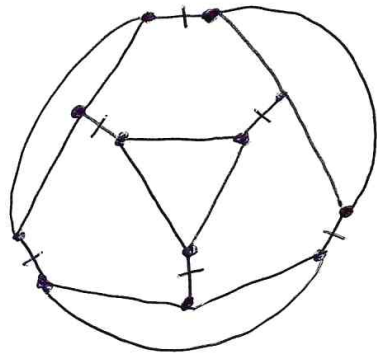


The four-color theorem asserts the colorability of planar isthmus-free trivalent graphs. From any given PM of a graph we can obtain other PM's by switching cycles consisting of alternating edges and PM edges. For example:



Problem  
 Understand relations among the links  $K(\mathbb{Q}, M)$  as  $M$  runs over PM's of a given graph  $\mathbb{Q}$ .



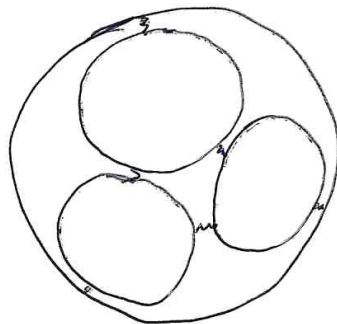


$(\mathcal{O}, M)$

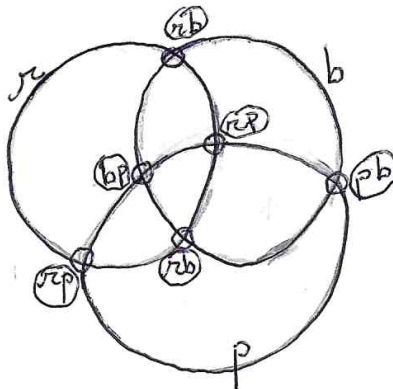
odd  
perfect  
matching

$$\{\cancel{X}\} = \{M\} + \{\cancel{X}\}$$

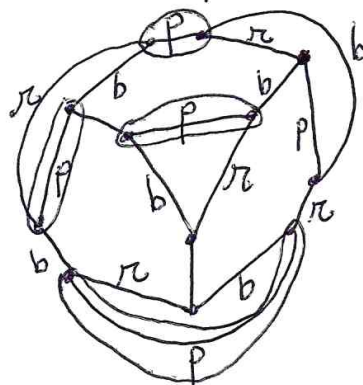
Basic Tautology  
applied with  
3 colors to any  
perfect matching.



$S$  state  
not colorable  
with 3 colors.



$S'$   
a  
3-colored  
state.



an even  
perfect  
matching



Transfer Basic Tautology to Virtual Link Diagrams:

$$\{ \times \} = \{ \cup \} + \{ \cap \}$$

$$( \times \longrightarrow \cup )$$

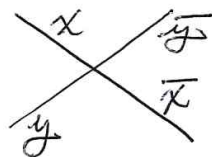
$$\{ 0 \} = 3.$$

THIS WILL LOCATE ALL POSSIBLE 3-COLORINGS, BUT IT IS NOT INVARIANT UNDER REIDEMEISTER MOVES.

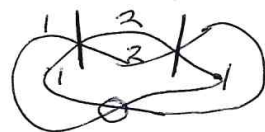
BUT: We can restrict to 2 colors.

This corresponds to counting all possible colorings for a perfect matching where all matching edges have a constant color (e.g. ③).

$$[K] = \sum_{\text{all 2-colorings of } K} 1$$



e.g.



$$\begin{aligned} 2 &= \overline{1} \\ 1 &= \overline{2} \end{aligned}$$

And generalize to

$$[K] = \sum_{\sigma \in 2\text{-colorings of } K} \langle K | \sigma \rangle$$

$$\langle \text{X} | \text{m} \rangle = A$$

$$\langle \text{X} | \text{z} \rangle = A^{-1}$$

Binary Bracket

$$\boxed{\begin{aligned} [\text{X}] &= A[\text{m}] + A^{-1}[\text{z}] \\ [0] &= 2 \end{aligned}}$$

$$\Rightarrow \begin{cases} [K] \text{ is invariant under } R2 \text{ and } R3. \\ [\text{m}] = A[\sim] \\ [\sim] = A^{-1}[\text{m}] \end{cases}$$

$$\underline{At A=1: [K] \neq \emptyset \iff}$$

When  $A=1$   
 $[K] \neq 0$

$$\Rightarrow |[K]| = 2$$

$K$  corresponds to an even PM.  
 # components of  $K$   
 # cycles in  $(G-M)$   
 $(G = K^{-1}(K))$ .

Thus the binary bracket <sup>(A=1)</sup> and the standard bracket <sup>(A=i)</sup> both detect even PM'S. The binary bracket and the standard bracket (aka Jones polynomial) both contain very definite topological information about PM'S.

And the binary bracket is a specialization of the Basic Tautology that handles all possible 3 colorings of a trivalent graph.

Can we use this relationship of topology and combinatorics to unlock the coloring problem for trivalent graphs?

### 6.° And Khovanov Homology

⑤ a trivalent graph with PM, M.  
Define  $Kho^*(G, M) = Kho^*(K(G, M))$   
where  $Kho^*$  is your favorite version of Khovanov Homology for virtuals. We are presently working with William Rushworth's Doubled Khovanov Homology and relating the Lee Homology part of that to the 2-coloring of virtual link diagrams.

arXiv:  
1704.0734

