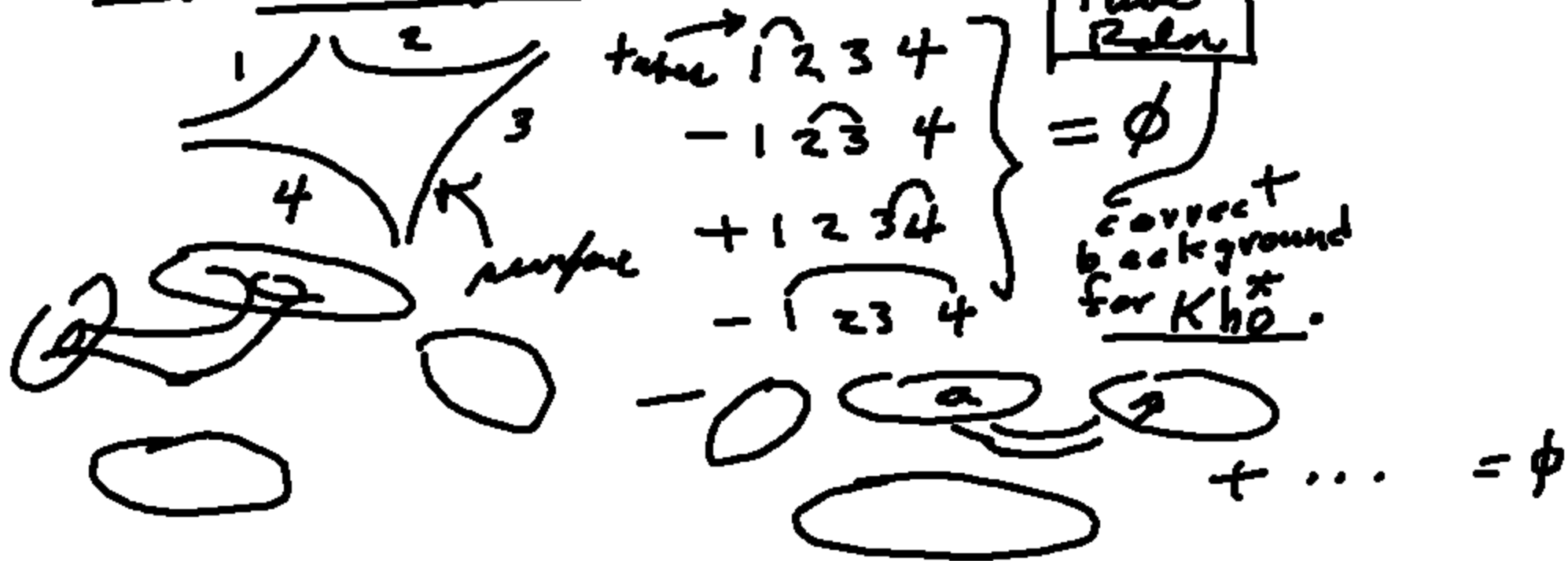


- One more lecture in the course next week.
- A problem set will be distributed in few days.
- We may continue in January some form.

1. Tubing Relation



$$\partial(A, B, C, D) = (\overbrace{A}^{\leftarrow \text{a "product" }}, B, C, D) \\ - (A, \overbrace{BC}, D) \\ + (A, B, \overbrace{CD}) \\ - (\overbrace{A, B, C}, D)$$

$$\partial(A, B, C, D) = (AB, C, D) \\ - (A, BC, D) \\ + (A, B, CD) \\ - (DA, B, C)$$

What is needed
for $\partial \circ \partial = \phi$?

$$\partial\partial(A, B, C, D)$$

$$= \partial \left[(AB, C, D) - (A, BC, D) + (A, B, CD) - (DA, B, C) \right]$$

$$= \begin{aligned} & \underline{(AB)C, D} - \underline{A(BC), D} + \underline{AB, CD} \\ & - \underline{(AB, CD)} + \underline{A, (BC)D} - \underline{(A, B)CD} \\ & + \underline{D(AB), C} - \underline{(DA, BC)} + \underline{(CD)A, B} \\ & \qquad \qquad \qquad - \underline{(DA)B, C} \\ & \qquad \qquad \qquad + \underline{(DA, BC)} \\ & \qquad \qquad \qquad - \underline{(C(DA), B)} \end{aligned}$$

$$= \emptyset \quad \text{exactly when op. } A, B \rightarrow AB \text{ assoc.}$$



Tube Relation looks like
 setting $\partial \alpha = \emptyset$ in a
 complex made from cobordism.

But at this stage the
 boundary map is not
 quite alg defined.

"Cobs / ∂ Cobs"

What will come of
 further thought
 along this line?

The 4-Tube Relation is somehow
 analogous to the (famous)
4-Term Relation for Vassiliev
Invariants for Knots and Links.

An invar I of knots / links is
 said to be a Vassiliev Invariant

$$I \nearrow \nearrow - I \nearrow \searrow = I \times \nearrow$$

singular crossings

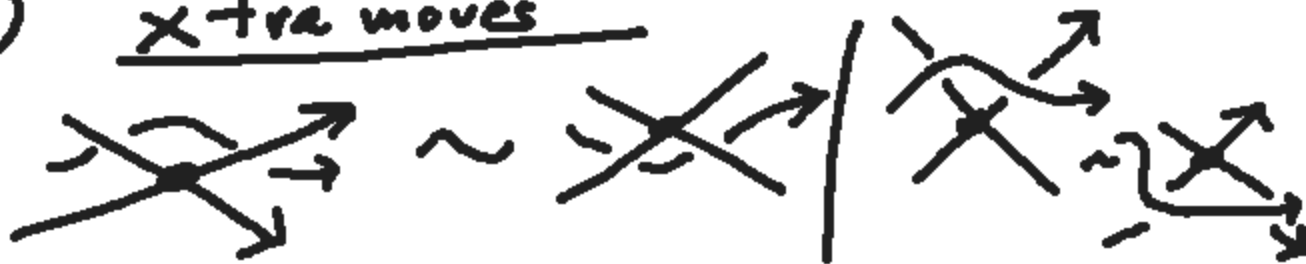
The invar is defined on
 knots with sing crossings

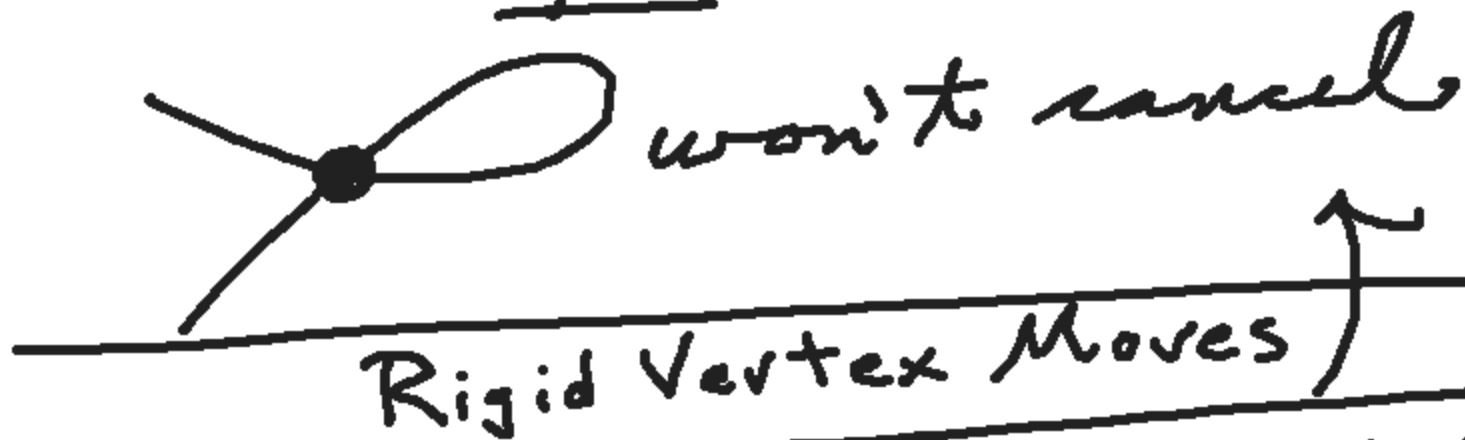
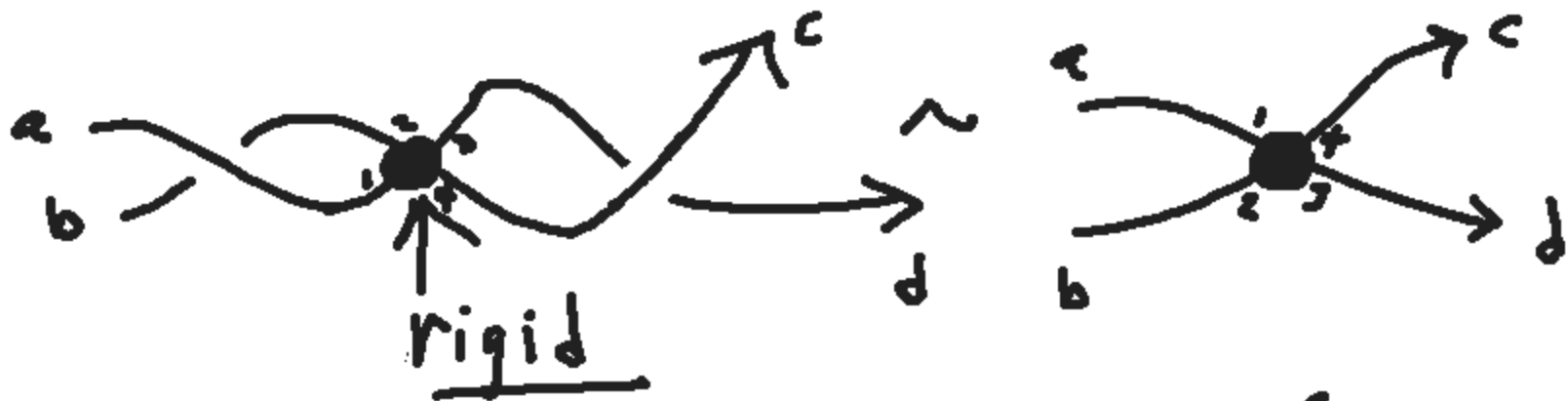


Graphs
 $\subset \mathbb{R}^3$



xtra moves





Fact. If J_K is an invariant of knots
 + links then we extend
 J to RVGraphs via J_{extended} is
 an RVG inv.

$$J_{\text{cross}} = a J_{\text{smooth}} + b J_{\text{other}}$$

eval J on the knots + links.

e.g. Let $J_K(q) = \text{Jones poly.}$

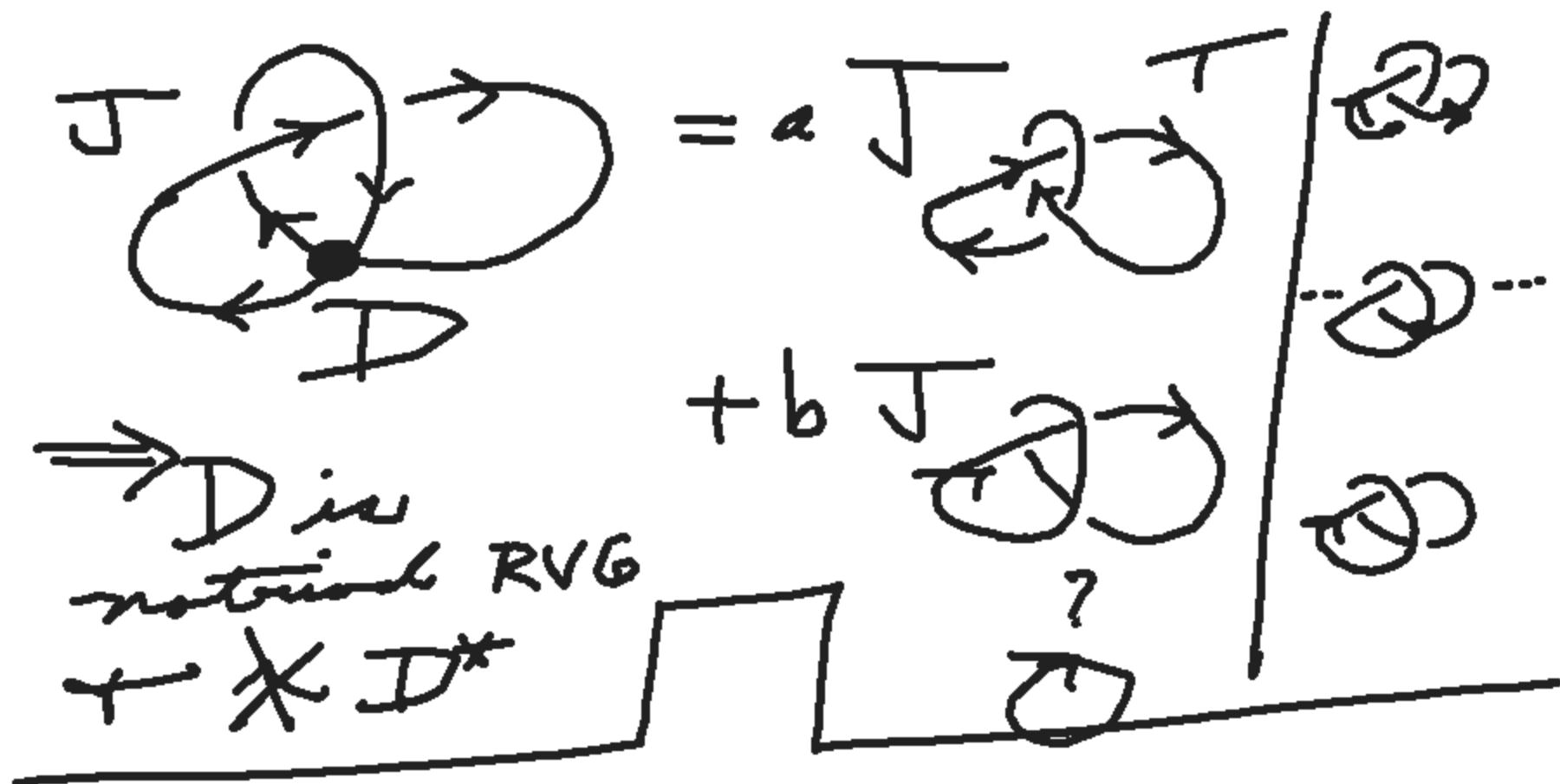
$$J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \bullet \\ \diagup \quad \diagdown \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \stackrel{=}{=} a J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} + b J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \diagdown \quad \diagup \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array}$$

$$= a J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} + b J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \diagdown \quad \diagup \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array}$$

$$\stackrel{=}{=} J \begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \text{---} \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \quad \checkmark$$




$$J \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \stackrel{=}{=} a J \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} + b J \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array}$$

$$= a J \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} + b J \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = J \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array}$$






- $I \xrightarrow{\rightarrow} - I \xrightarrow{\rightarrow} = I \xrightarrow{\times}$
- I is said to be of finite type N
 If $I \xrightarrow{\times} \dots \xrightarrow{\times} = \emptyset$. Vanishes for G with $> N$ nodes.

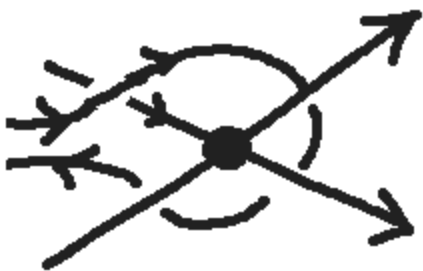
Lemma. If I is Vasiliev
 invariant of type N then I_G
 is independent of $G \subset \mathbb{R}^3$ for
 G with exactly N nodes.

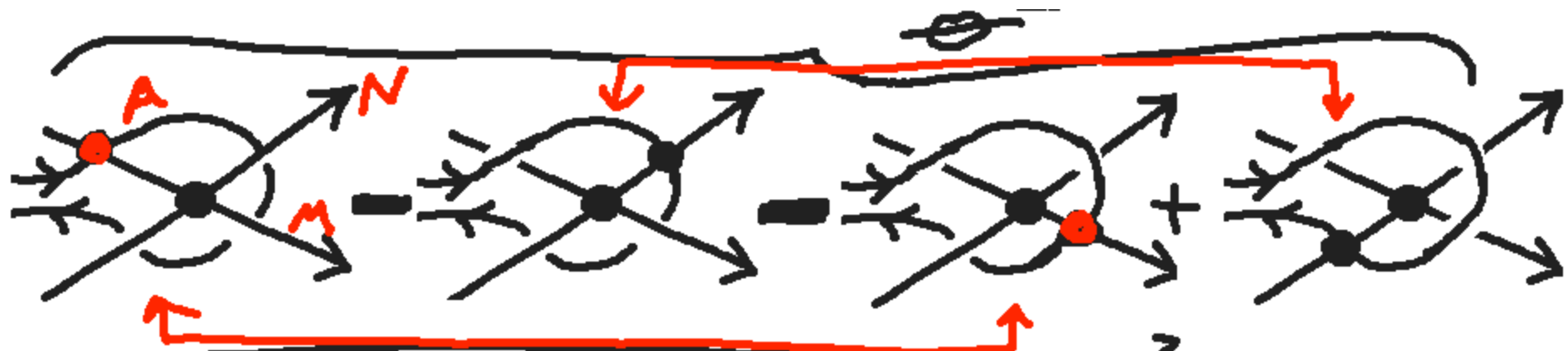
$N=2$ I  $-\frac{I}{\phi}$  $= I$ 

Could
 assume
 $I_{\bigcirc} = \phi$
 $I_{\text{cross}} = \phi$
 " $I_{\bigcirc} - I_{\text{cross}}$

I  $-\frac{I}{\phi}$  $\xrightarrow[\text{same } I]{}$

$= I$  $= \phi$
Type 2

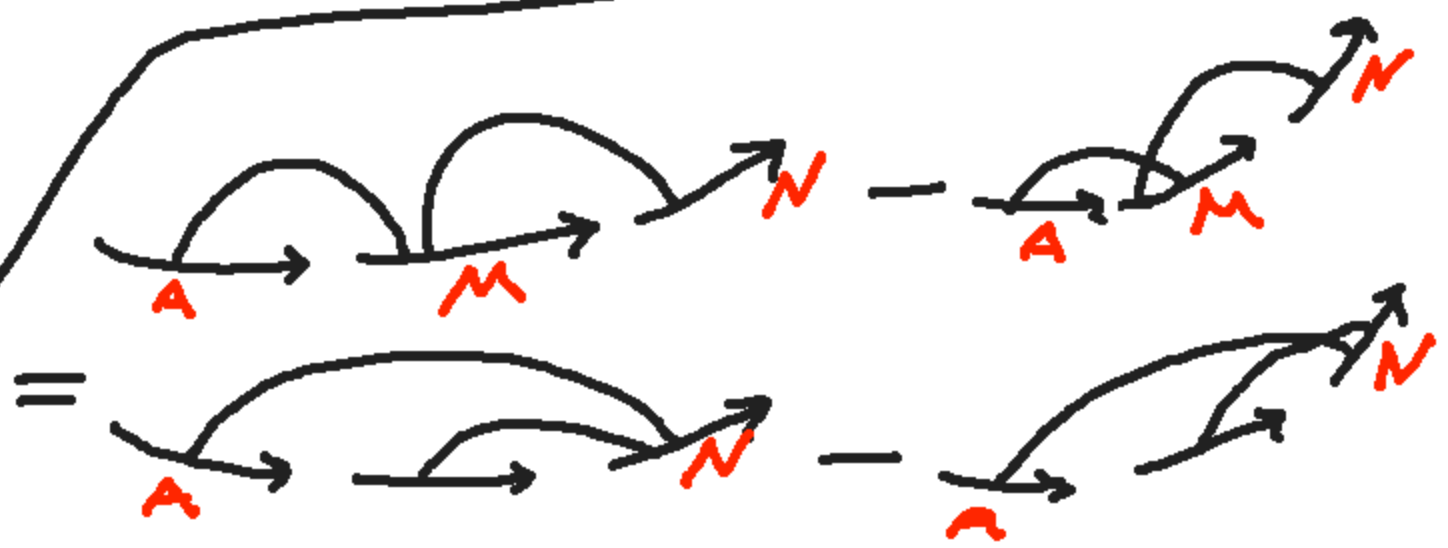


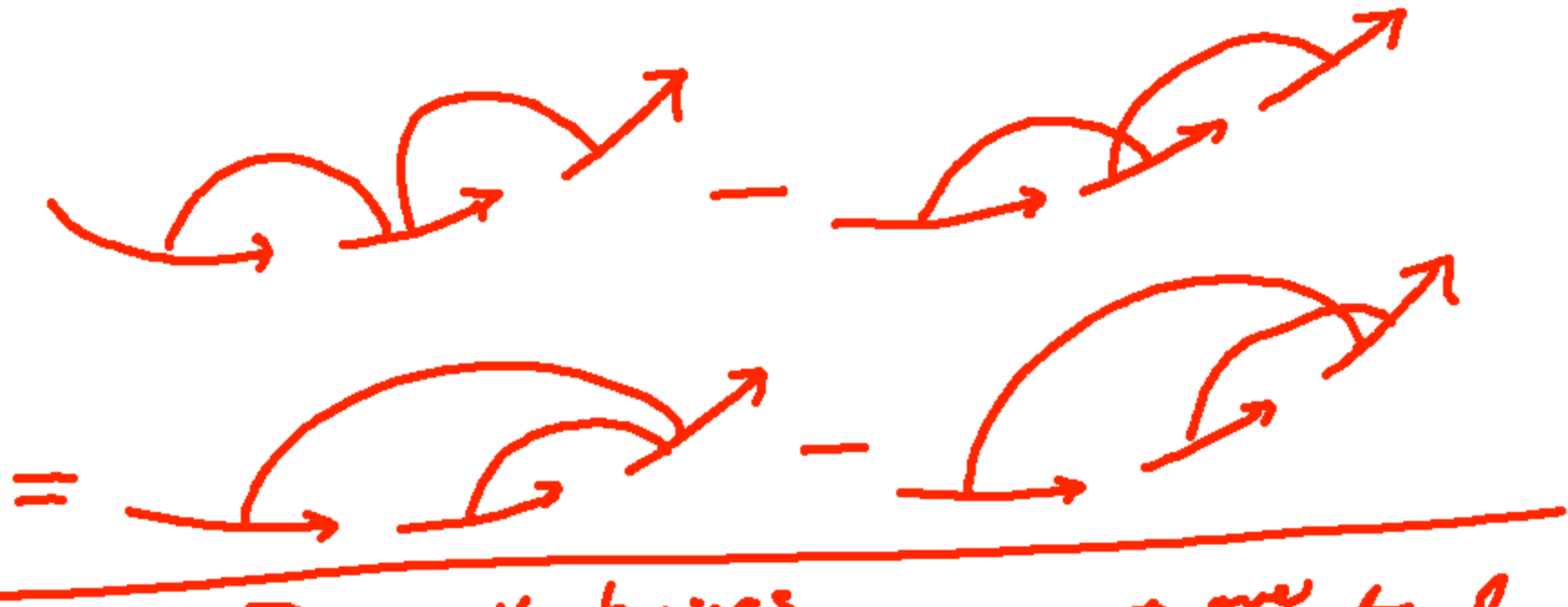


1212
 Gauss
 code for
nodes



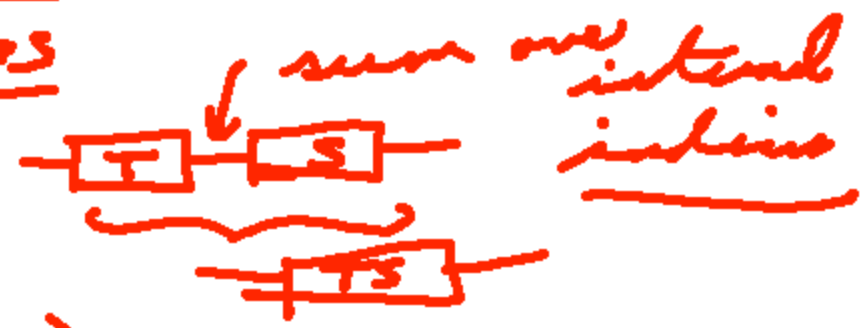
on Gauss
 Diags.





Diag Matrices

$$\sum_j T_{ij} \delta_{jk} = (TS)_{ik}$$



$$\sum_j T_{ij} \delta_{jk} = (TS)_{ik}$$

Lie Alg $\{T^a\}$: $T^a T^b - T^b T^a \Rightarrow \sum_i f_i^{ab} T^i$
 basis $a=1, \dots, d$ $[T^a, T^b] \in \text{alg spanned by } \{T^c\}$

$$T^a T^b = T^b T^a + f_c^{ab} T^c$$

Let a, b be any two generators in \mathfrak{g}

$$[T^a, T^b] = f_c^{ab} T^c$$

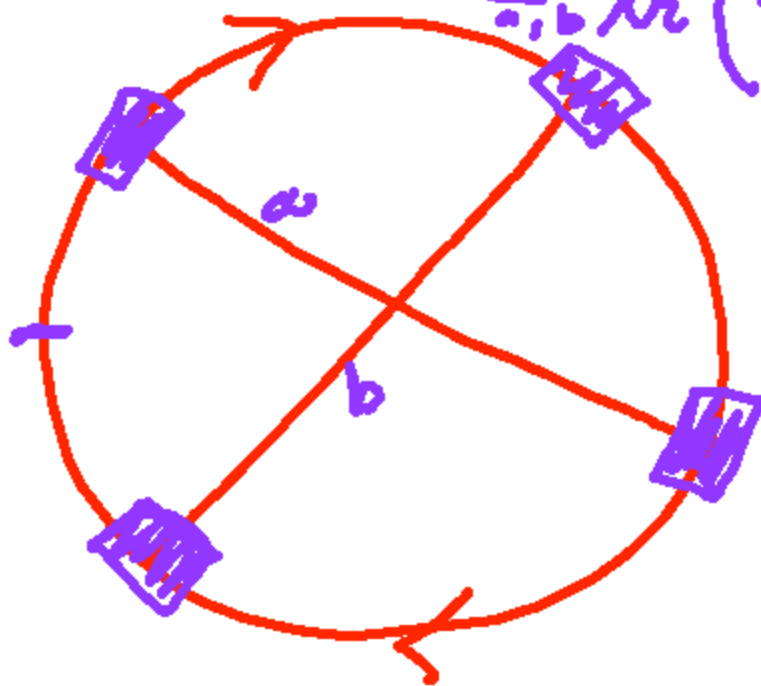


Lie Alg
↔
4TR



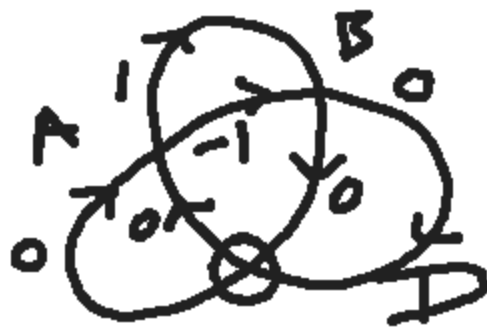
$$\sum_{a,b} \text{tr}(T_a T_b T_a T_b)$$

||
eval of



Dissection on affine braid

$Wt(D)$



	w_+	w_-
A	-1	1
B	1	-1



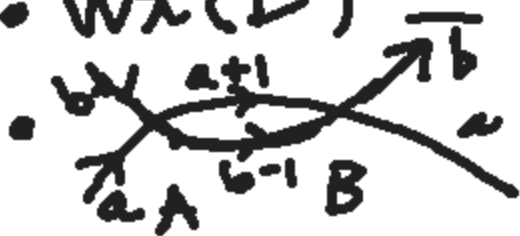
$$P_K = \sum_c \text{sgn}(c) (t^{w_{\text{sgn}(c)}} - 1)$$

$$P_K = t^{-1} + t - 2$$

$Wt(D)$

Wt F at Viatches

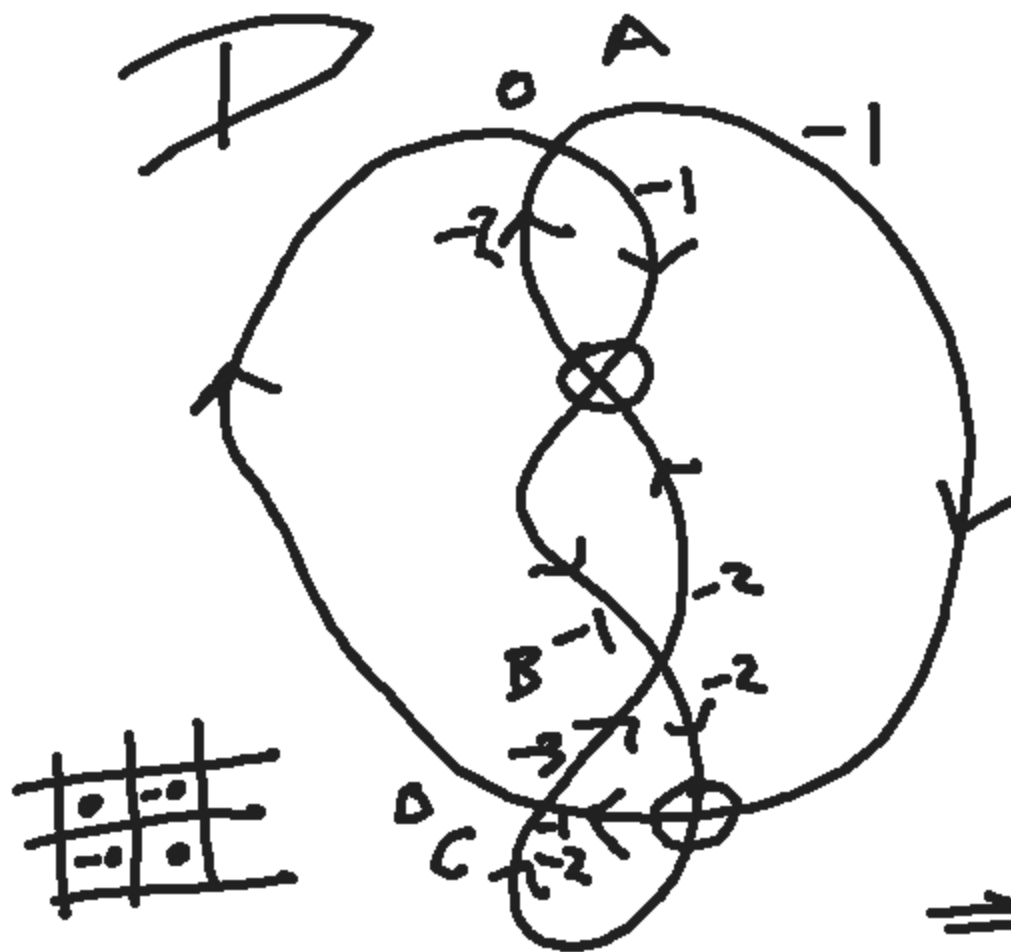
• $Wt(D) = Wt(D')$ if $D \sim D'$



	w_+	w_-
A	•	*
B	*	•

$D \sim D'$
RS

A	$b-a-1$	$a-b+1$
B	$a+1-b$	$b-1-a$



	w_+	w_-
A	1	-1
B	1	-1
C	-2	2

} $W(D)$

Fact: D, D^r
 $D^r = D$ with reversed orientation.

$\Rightarrow \text{Wt}(D^r) = \text{Column reversal of } \text{Wt}(D)$

Let $\hat{W}(D) = \text{reversal of } W(D)$

R3. $W(D) \neq \hat{W}(D)$ preserved.

R2. $W(D) \neq \hat{W}(D)$ preserved

} If $W(D) \neq \hat{W}(D)$
 $\Rightarrow D$ not equiv
to D^r .